

one second, then making the user wait 60 seconds before he can do any more searches.

The algorithm can be further modified to include a higher maximum capacity, allowing up to 10 tokens to be stored, for example, with a rate limit of five per minute, allowing users to initially do 10 searches in the first minute but only five per minute afterward (until they give the system a minute or two to recover).

This limiting can be done on a global basis, a per-IP address basis (simply have an array list of IPs and their current allowed searches and last time), or even a per-user basis (i.e., one leaky bucket per user), or the methods can be combined, allowing each user five searches per minute and a system total of 50 searches per minute (assuming more would make your system become too slow to use).

Distributed Token Buckets

A common flaw made with rate-limiting systems that occurs in applications with more than one server (which is almost all major applications nowadays) is that the various servers and system components do not properly share state. For example, a site with five front-end web servers may implement a download rate limit or a limit on login attempts for clients, but if they do not communicate

with each other, an attacker would be able to execute five times the rate-limited downloads or attacks by hitting each server at the same time. Of course, the solution is to use a shared state. Two possible solutions are to use a database with row-level locking (reducing the amount of contention for checking and updating the shared state information) or use something very lightweight that is also extremely fast, such as memcached [4]. The memcached server not only supports storing keys and values associated with the keys but can also store a counter (a 64-bit integer) that can be incremented or decremented atomically (i.e., multiple processes won't step on each others' toes) with the *incr* and *decr* commands. In this way, you can simply insert a key with a value equivalent to the user name, IP address, or whatever you are using and then maintain a counter. For each attempt to log in, you increment the counter, and you regularly decrement the counter by a control process to ensure that attempts are expired.

Lockouts vs. Timeouts

Another aspect of rate limiting is the ability to reduce problems caused by lockouts – systems that try to protect themselves by locking out users after a certain number of failures, for example. Lockouts are prone to spoofing attacks.

If attackers can make themselves look like legitimate users (i.e., by trying to use the victim's username to log in), they might be able to lock out that user (assuming you lock an account out after three bad passwords). Alternatively, if you have customers or users coming from behind proxy servers, you might end up with one bad user blocking access for a group of legitimate users.

The use of increasingly long timeouts (such as doubling the waiting period each attempt) can be just as effective as a lockout but has the added (dis)advantage that when the attack stops, the timeout will eventually reset, allowing whatever caused the timeout to start again. Of course, depending on what you want, you can set the timeouts appropriately to either allow users into their accounts or block spammers for extended periods of time.

What Is Your Quest

A final option available in conjunction with rate limiting is allowing users of a system to prove that they do not have hostile intentions (despite their "bad" behavior). For example, if you send too many similar-looking queries to Google in a short time period you might get a web page saying "Sorry" that directs you to enter a CAPTCHA string to prove you are human and not an automated program or malicious program. All of which is much better than a user staring at a blank and non-responsive web page. ■

Listing 1: Leaky Bucket Pseudo-Code

```
01 searches=5
02 per_second=60
03 current_allowed=searches
04 last_check = time()
05 while process(search_terms):
06     # we determine how many seconds since the last search
07     time_now=time()
08     time_passed = time_now - checked_at
09     # and set when our last search was
10     checked_at=time_now
11     # add tokens to bucket:
12     current_allowed += time_passed * (searches / per_second)
13     # check if we have any tokens
14     if (current_allowed > searches):
15         # we have reached our max. tokens
16         current_allowed = searches
17     if (current_allowed < 1):
18         #partial token, ignore search
19         discard_search()
20     else:
21         # we have at least one token
22         do_search()
23         # and we "spend" the token
24         current_allowed = current_allowed - 1
```

INFO

- [1] "Apache HTTPD" by Kurt Seifried, *Linux Magazine*, September 2009, pg. 52, <http://www.linuxpromagazine.com/Issues/2009/106/APACHE-HTTPD>
- [2] Leaky Bucket: http://en.wikipedia.org/wiki/Leaky_bucket
- [3] Token Bucket: http://en.wikipedia.org/wiki/Token_bucket
- [4] Memcached: <http://memcached.org/>

THE AUTHOR

Kurt Seifried is an Information Security Consultant specializing in Linux and networks since 1996. He often wonders how it is that technology works on a large scale but often fails on a small scale.

